

PATENT PD-200095A

**APPLICATION OF COMPLEX CODES TO MAXIMIZE USER LINK
UTILIZATION**

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UTILIZATION

BACKGROUND OF THE INVENTION

0945433-033104
The present invention relates generally to code
division multiple access (CDMA) cellular communications
systems. More specifically, but without limitation
thereto, the present invention relates to assigning
5 variable length codes to service a maximum number of users
in a code division multiple access service area that is
assigned a fixed frequency band.

The radio frequency spectrum is limited almost
everywhere in the world and is generally licensed in fixed
10 frequency bands. Code division multiple access (CDMA)
cellular communications systems typically use a set of
CDMA codes that are orthogonal to one another to avoid
mutual interference. The term orthogonal is applied to a
set of codes if the vector dot product of any code in the
15 set with any other code in the set results in zero. For
example, the code (-1, -1) is orthogonal to the code (-1,
1) since $-1 \cdot -1 + -1 \cdot 1 = 0$. Currently known orthogonal
CDMA codes have a code length of 2^n , where n is a positive
integer. The available frequency bands are not
20 necessarily integrally divisible by 2^n , however. For
example, for a user link that has 5 MHz of available
bandwidth, the coded bit rate is 144 kbps (kilobits per
second). If quadrature phase shift keying (QPSK)
modulation is used, then the maximum ratio of bandwidth to

symbol rate, or spreading code length, would be 68. However, no orthogonal codes are available having a code length of 68. The closest power of two that is less than 68 is 64. About six percent of the available bandwidth is therefore unusable using conventional CDMA codes. In other cases, there may be a substantially higher percentage of unused bandwidth due to the limited code lengths available using conventional CDMA codes. The unused bandwidth represents a loss of revenue from potential CDMA subscribers that might otherwise be included in the same service area.

SUMMARY OF THE INVENTION

The present invention advantageously addresses the problems above as well as other problems by providing a method of increasing utilization of user link bandwidth for a code division multiple access communications system.

In one embodiment, the present invention may be characterized a method of increasing utilization of user link bandwidth for a code division multiple access communications system that includes the steps of selecting a set of orthogonal complex codes each having a code length that is greater than a code length of an optimum real code and less than or equal to a spreading code length; and transferring symbols across at least one of a plurality of user links to or from at least one of a corresponding plurality of user terminals wherein the symbols are represented by a corresponding one of the set of orthogonal complex codes.

In another embodiment, the present invention may be characterized as a code division multiple access

communications system that includes a base station; a geo-
stationary platform; a feeder link coupled to the base
station and the geo-stationary platform for transferring
symbols between the base station and the geo-stationary
5 platform; a plurality of user terminals; and a plurality
of user links coupled respectively to the plurality of
user terminals and to the geo-stationary platform for
transferring symbols between the geo-stationary platform
and at least one of the plurality of user terminals
10 wherein the symbols are represented by at a corresponding
one of a set of orthogonal complex codes having a code
length that is greater than a code length of an optimum
real code and less than or equal to a spreading code
length.

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BRIEF DESCRIPTION OF THE DRAWINGS

The features and advantages of the present
invention may be apprehended from the following
description thereof, presented in conjunction with the
20 following drawings wherein:

FIG. 1 is a diagram of a code division multiple
access (CDMA) communications system of the prior art; and

FIG. 2 is an exemplary set of orthogonal complex
codes for the maximizing the user link bandwidth of the
25 code division multiple access communications system of
FIG. 1 according to an embodiment of the present
invention.

Corresponding reference characters indicate
corresponding elements throughout the several views of the
30 drawings.

DESCRIPTION OF THE EMBODIMENTS

FIG. 1 is a diagram of a typical code division multiple access (CDMA) communications system 100. Shown in FIG. 1 are a base station 102, a feeder link 104, a
5 geo-stationary platform 106, user links 108, and user terminals 110.

The base station 102 transmits and receives symbols between user terminals 110 across the feeder link 104. The feeder link 104 is an RF communications link
10 between the base station 102 and the geo-stationary platform 106. The geo-stationary platform 106 is typically a geo-stationary satellite, however other platforms having a relatively fixed position with respect to the base station 102 may also be used according to
15 techniques well known in the art. Because the feeder link 104 is between only two points (point-to-point), code division multiple access communications between the base station 102 and the geo-stationary satellite 106 are generally synchronous, and may be modulated by, for
20 example, quadrature phase shift keying (QPSK).

The user terminals 110, may be, for example, cellular telephones. The user links 108 are RF communications links between the respective user terminals 110 and the geo-stationary satellite 106. Because the
25 user links 108 are between multiple points whose positions relative to the geo-stationary satellite 106 may change (point-to-multiple-point), communications between the user terminals 110 and the geo-stationary satellite 106 are generally asynchronous, and may be modulated, for example,
30 by quadrature phase shift keying (QPSK).

To avoid mutual interference, code division multiple access (CDMA) communications systems generally use mutually orthogonal codes for transferring symbols across the user links 108. One example of mutually
5 orthogonal codes is a set of Walsh codes. Walsh codes may be generated from Walsh code functions as follows. To generate a Walsh code sequence of code length two starting from a seed value of -1 for a Walsh code of code length one, the seed value of -1 is appended to itself to
10 generate the first Walsh code, (-1,-1). The bit-inverse of the seed value is appended to the seed value to generate the second Walsh code, (-1,1). This is the equivalent of creating the 2-bit Walsh code table shown in Table 1 below.

15

| Walsh Code | W0 | W1 |
|------------|----|----|
| 0 | -1 | -1 |
| 1 | -1 | 1 |

TABLE 1

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The procedure used to create Table 1 may be repeated to generate the Walsh codes for the next higher order, i.e., the next longer Walsh code length, as shown in Table 2 below. The 2 X 2 matrix in bit positions W0
25 and W1, also called the Hadamard transform matrix, is appended to itself in bit positions W2 and W3 to generate the first two Walsh codes 0 and 1. The second two Walsh codes 2 and 3 are generated by appending the bit-inverse of the 2 X 2 matrix to the original 2 X 2 matrix.

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| Walsh Code | W0 | W1 | W2 | W3 |
|------------|----|----|----|----|
| 0 | -1 | -1 | -1 | -1 |
| 1 | -1 | 1 | -1 | 1 |
| 2 | -1 | -1 | 1 | 1 |
| 3 | -1 | 1 | 1 | -1 |

TABLE 2

Additional Walsh code tables for the higher order Walsh code lengths may be generated by repeating the procedure above. For example, to create eight-bit Walsh codes, the 4 X 4 matrix of Table 2 is replicated three times and inverted in the lower right hand quadrant as shown in Table 3 below.

15

| Walsh Code | W0 | W1 | W2 | W3 | W4 | W5 | W6 | W7 |
|------------|----|----|----|----|----|----|----|----|
| 0 | -1 | -1 | -1 | -1 | -1 | -1 | -1 | -1 |
| 1 | -1 | 1 | -1 | 1 | -1 | 1 | -1 | 1 |
| 2 | -1 | -1 | 1 | 1 | -1 | -1 | 1 | 1 |
| 3 | -1 | 1 | 1 | -1 | -1 | 1 | 1 | -1 |
| 4 | -1 | -1 | -1 | -1 | 1 | 1 | 1 | 1 |
| 5 | -1 | 1 | -1 | 1 | 1 | -1 | 1 | -1 |
| 6 | -1 | -1 | 1 | 1 | 1 | 1 | -1 | -1 |
| 7 | -1 | 1 | 1 | -1 | 1 | -1 | -1 | 1 |

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TABLE 3

Each Walsh code has a code length of 2^n bits, where n is a power of 2. In quadrature phase shift keying modulation, a Walsh code is represented by modulating the RF carrier phase between -90 degrees for a -1 and +90 degrees for a 1. Because -90 degrees and +90 degrees lie

on the real axis of a phase vector coordinate system, the corresponding codes are called real codes.

In a code division multiple access system wherein orthogonal codes are generated by Walsh functions, the user data rate may be variable. For example, for a set of orthogonal Walsh codes having a code length of eight and a basic data rate of R , a code division multiple access system can support the following scenarios:

- (1) all eight users with a data rate of R ,
- (2) one user with a data rate of $2R$ and 4 users with a data rate of R ,
- (3) one user with a data rate of $4R$ and 2 users with a data rate of R , and
- (4) one user with a data rate of $8R$. In this case, the entire bandwidth is dedicated to a single user.

Other scenarios such as those described above may be used to show that the number of available data rates for a given code length is half the code length. In the example above, there are four data rates (R , $2R$, $4R$, and $8R$).

By introducing additional phase values for modulating the phase of the RF carrier, a set of orthogonal complex codes to represent symbols transferred across the user links may be selected. Each of the orthogonal complex codes has a selected code length other than 2^n . For example, a selected orthogonal complex code may have a code length of $4n$. Because there are more values of $4n$ in a given numerical interval than there are values of 2^n , a set of orthogonal complex codes may be selected to utilize a greater portion of a fixed bandwidth than may be possible with the optimum set of real codes

having a code length equal to the closest value of 2^n that does not exceed the spreading code length.

By way of example, a set of orthogonal complex codes of code length $4n$ to represent symbols transferred across the user links 108 may be generated from the following Kronecker tensor product:

$$C_{L \times P} = A_L \otimes W_P \quad (1)$$

10 wherein

$C_{L \times P}$ is a matrix of complex codes each having a code length equal to $L \times P$,

L is a positive integer,

P equals 2^n and n equals a positive integer,

15 W_P is a Walsh code of code length P ,

A_L is a coefficient matrix of elements a_{jk} , where $j = 1 \dots L$, $k = 1 \dots L$, and

$$a_{jk} = e^{j2\pi(j-1)(k-1)/L} \quad (2)$$

20

The optimum Walsh code rate is the highest power of two that is less than or equal to the ratio of bandwidth to symbol rate, or spreading code length. For example, the optimum Walsh code length for a spreading code length of 12 is 2^3 , since the next higher Walsh code length of 2^4 exceeds the spreading code length. Because orthogonal complex codes generated, for example, by formula (1) may have a code length of $4n$, a set of orthogonal complex codes having a code length greater than the optimum Walsh code length may be generated to fully utilize a spreading code length of 12. For a spreading

code length of 12, the following mixture of data rates is possible:

- (1) all 12 users with a data rate of R ,
- (2) one user with a data rate of $3R$ and nine users with a data rate of R ,
- (3) one user with a data rate of $6R$ and six users with a data rate of R , and
- (4) one user with a data rate of $12R$. In this case, the entire bandwidth is dedicated to a single user.

By way of example, to generate a set of complex codes of code length 12 using formula (1), let $L = 3$ and $P = 4$. The complex CDMA code is then given by the following relation:

$$C_{3 \times 4} = \begin{bmatrix} a_{11}W_4 & a_{12}W_4 & a_{13}W_4 \\ a_{21}W_4 & a_{22}W_4 & a_{23}W_4 \\ a_{31}W_4 & a_{32}W_4 & a_{33}W_4 \end{bmatrix} \quad (3)$$

Numerically, $C_{3 \times 4}$ equals the matrix of orthogonal complex codes illustrated in FIG. 2. Each column and corresponding row of the matrix of orthogonal complex codes illustrated in FIG. 2 is a complex code that is orthogonal to each of the other complex codes.

For the example of a given bandwidth of 5 MHz and a given bit rate of 144 kbps described above, the bandwidth utilization of the user link in the code division multiple access communication system 100 may be improved from $64/68 = 94$ percent for the optimum set of real codes to $68/68 = 100$ percent by implementing a set of orthogonal complex codes such as those generated in the example above as well as other sets of orthogonal complex

codes that may be generated for specific applications.
The increase in bandwidth utilization is generally
available in those cases where a set of orthogonal complex
codes may be generated having a code length that is
5 greater than the optimum real code length and less than or
equal to the corresponding spreading code length for a
given user link bandwidth and a given data bit rate. The
greater code length of the selected set of orthogonal
complex codes allows more user terminals 110 to share the
10 same service area, thereby reducing the number of service
areas required and the corresponding total cost of support
equipment. In cases where the ratio of the spectrum
bandwidth to the symbol rate happens to be equal or very
close to 2^n , real codes are preferable to complex codes.
15 At other ratios, however, complex codes may be used
advantageously to increase the capacity of a code division
multiple access communications system.

Other modifications, variations, and
arrangements of the present invention may be made in
20 accordance with the above teachings other than as
specifically described to practice the invention within
the spirit and scope of the following claims.